#### Scalable Tools for Debugging Non-Deterministic MPI Applications \_\_\_\_\_\_ ReMPI: MPI Record-and-Replay tool \_\_\_\_\_

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# Debugging large-scale applications is already challenging



[1] Source: http://www.prweb.com/releases/2013/1/prweb10298185.htm, CAMBRIDGE, UK (PRWEB) JANUARY 08, 2013

With trends towards asynchronous communication patterns in MPI applications, <u>MPI non-determinism</u> will significantly increase debugging cost





#### What is MPI non-determinism ?

- Message receive orders can be different across executions
  - Unpredictable system noise (e.g. network, system daemon & OS jitter)
- Floating point arithmetic orders can also change across executions



Execution A: (a+b)+c



Execution B: a+(b+c)



#### Non-determinism also increases debugging cost

Control flows of an application can change across different runs



- Non-deterministic control flow
  - Successful run, seg-fault or hang
- Non-deterministic numerical results
  - Floating-point arithmetic is non-associative

 $(a+b)+c \neq a+(b+c)$ 

➔ Developers need to do debug runs until the target bug manifests

In non-deterministic applications, it's hard to reproduce bugs and incorrect results. It costs excessive amounts of time for "<u>reproducing</u>" target bugs





#### Non-deterministic bugs ---- Case study: Pf3d and Diablo/Hypre 2.10.1

 Debugging non-deterministic hangs often cost computational scientists substantial time and efforts



- Diablo hung only once every 30 runs after a few hours
- The scientists spent 2 months in the period of 18 months and gave up debugging it

Hypre is an MPI-based library for solving large, sparse linear systems of equations on massively parallel computers



- Pf3d hung only when scaling to half a million MPI processes
- The scientists refused to debug for 6 months ...



#### **Non-deterministic numerical result**

--- Case study: "Monte Carlo Simulation" (MCB)

CORAL proxy applicationMPI non-determinism

#### Table 1: Catalyst Specification

¥	
Nodes	304 batch nodes
CPU	2.4  GHz Intel Xeon E5-2695 v2
	(24  cores in total)
Memory	128 GB
Interconnect	InfiniBand QDR (QLogic)
Local Storage	Intel SSD 910 Series
	(PCIe 2.0, MLC)



#### Final numerical results are different between 1<sup>st</sup> and 2<sup>nd</sup> run



\* The source was modified by the scientist to demonstrate the issue in the field



## Why MPI non-determinism occurs ?

- It's typically due to communication with MPI\_ANY\_SOURCE
- In non-deterministic applications, each process doesn't know which rank will send message
- Messages can arrive in any order from neighbors  $\rightarrow$  inconsistent message arrivals





#### **ReMPI can reproduce message matching**

ReMPI can reproduce message matching by using record-and-replay technique



- Traces, records message receive orders in a run, and replays the orders in successive runs for debugging
  - Record-and-replay can reproduce a target control flow
  - Developers can focus on debugging a particular control flow





#### **Record overhead to performance**

Performance metric: how many particles are tracked per second



- ReMPI becomes scalable by recording to local memory/storage
  - Each rank independently writes record  $\rightarrow$  No communication across MPI ranks







#### **Record-and-replay won't work at scale**

- Record-and-replay produces large amount of recording data
  - Over "<u>10 GB/node</u>" per day in MCB
  - Over "<u>24 GB/node</u>" per day in Diablo
- For scalable record-replay with low overhead, the record data must fit into local memory, but capacity is limited
  - Storing in shared/parallel file system is not scalable approach
  - Some systems may not have fast local storage



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#### **Clock Delta Compression (CDC)**







#### Logical clock vs. wall clock

"The global order of messages exchanged among MPI processes are very similar to a logical-clock order (e.g., Lamport clock)"



Lamport clock values of received messages for particle exchanges in MCB (MPI rank = 0)



### **Clock Delta Compression (CDC)**

 Our approach, clock delta compression, only records the difference between received order and logical order instead of recording entire received order





# Logical clock order is reproducible [1]

 Logical-clock order is always reproducible, so CDC only records the permutation difference





**Theorem 1.** CDC can correctly replay message events, that is,  $\boldsymbol{E} = \hat{\boldsymbol{E}}$  where  $\boldsymbol{E}$  and  $\hat{\boldsymbol{E}}$  are ordered sets of events for a record and a replay mode.

PROOF (MATHEMATICAL INDUCTION). (i) Basis: Show the first send events are replayable, i.e.,  $\forall x \text{ s.t. } ``E_1^x \text{ is send}$ events"  $\Rightarrow$  " $E_1^x$  is replayable". As defined in Definition 7.(i)  $E_1^x$  is deterministic, that is,  $E_1^x$  is always replayed. In Figure 12,  $E_1^1$  is deterministic, that is, is always replayed. (ii) Inductive step for send events: Show send events are replayable if the all previous message events are replayed, i.e., " $\forall E \to E$  s.t. E is replayed, E is send event set"  $\Rightarrow$  "E is replayable". As defined in Definition 7.(ii), E is deterministic, that is, E is always replayed. (iii) Inductive step for receive events: Show receive events are replayable if the all previous message events are replayed, i.e., " $\forall E \to E$  s.t. **E** is replayed, E is receive event set"  $\Rightarrow$  "E is replayable". As proofed in Proposition 1, all message receives in E can be replayed by CDC. Therefore, all of the events can be replayed, i.e.,  $E = \hat{E}$ . (Mathematical induction processes are graphically shown in Figure 12.)

Theorem 2. CDC can replay piggyback clocks.

PROOF. As proved in Theorem 1, since CDC can replay all message events, send events and clock ticking are replayed. Thus, CDC can replay piggyback clock sends.

[1] Kento Sato, Dong H. Ahn, Ignacio Laguna, Gregory L. Lee and Martin Schulz, "Clock Delta Compression for Scalable Order-Replay of Non-Deterministic Parallel Applications", In Proceedings of the International Conference on High Performance Computing, Networking, Storage and Analysis 2015 (SC15), Austin, USA, Nov, 2015.



# **Clock Delta Compression (CDC)**

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#### Implementation

- We use PMPI wrapper
  - Tracing message receive order
  - Clock piggybacking
- Clock piggybacking [1]



- When sending an MPI message, the PMPI wrapper defines a new
  - MPI\_Datatype that combines message payload & clock



[1] M. Schulz, G. Bronevetsky, and B. R. Supinski. On the Performance of Transparent MPI Piggyback Messages. In Proceedings of the 15th European PVM/MPI Users' Group Meeting on Recent Advances in Parallel Virtual Machine and Message Passing Interface, pages 194–201, Berlin, Heidelberg, 2008. Springer-Verlag.



#### **Compression improvement in MCB**

Total compressed record sizes on MCB at 3,072 procs (12.3 sec)



gzip itself can reduce the original format by 8x

#### 5x more reduction

- For example, if 1GB of memory per node for record-and-replay ...
  - w/o compression: 2 hours
  - gzip: 19 hours
  - CDC: 4 days

High compression

Compressed size becomes 40x smaller than original size





#### **Summary**

- Non-determinism is a common issue in debugging MPI applications
- ReMPI can help to reproduce buggy MPI behaviors with minimum record size



